



Validated non collision prediction of multiple drones

AID Meeting

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March, 13, 2020

Safety of an evolving swarm of drones
in term of collision avoidance

Two types of obstacles:

- **static**: forbidden areas;
- **dynamic**: the robots themselves



Goal

Guarantee on the **non collision** of the swarm with the environment (static obstacles) and the other drones (dynamic obstacles).

A drone i from the swarm is modeled with a controlled dynamical system:

$$(\mathcal{S}_i) \begin{cases} \dot{\mathbf{x}}_i = f(\mathbf{x}_i, \mathbf{u}_i) \\ \mathbf{x}_i(0) = \mathbf{x}_{i,0} \end{cases}$$

From a given control \mathbf{u}_i and a given time horizon T on which the control is applied to the system, the goal is to prove that, for two drones i and j :

$$x_i(t) \neq x_j(t), \forall t \in [0, T] \quad (1)$$

Uncertainties make this constraint intractable to check in general.

Validated method

Use of outer approximations to guarantee the non collision.

An interval is denoted $[x] = [\underline{x}, \bar{x}]$ with $\underline{x} \leq \bar{x}$.

The set of intervals is denoted

$$\mathbb{IR} = \{[x] = [\underline{x}, \bar{x}] \mid \underline{x}, \bar{x} \in \mathbb{R}, \underline{x} \leq \bar{x}\}.$$

The Cartesian product of intervals $[x] \in \mathbb{IR}^n$ is a **box**.

Interval arithmetic

Evaluating an arithmetic expression with intervals leads to an outer-approximation of the set it defines.

To deal with interval functions, an **interval inclusion function** (or interval extension) of a function can be defined.

Examples:

- **natural extension**: replaces the operations on reals by their interval counterparts using interval arithmetic;
- **mean value extension**: linearizes the function around its mean value.

Definition (IVP-ode)

An IVP-ODE is defined as

$$(\mathcal{S}) \begin{cases} \dot{\mathbf{y}} = f(t, \mathbf{y}) \\ \mathbf{y}(0) \in \mathcal{Y}_0 \subseteq \mathbb{R}^n, t \in [0, t_{\text{end}}] . \end{cases}$$

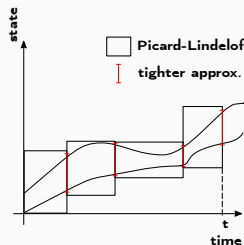
Goal is to compute $\mathbf{y}(t; \mathcal{Y}_0) = \{\mathbf{y}(t; \mathbf{y}_0) \mid \mathbf{y}_0 \in \mathcal{Y}_0\}$.

Phase 1 a priori enclosure $[\mathbf{y}_{[t_i, t_{i+1}]}]$ of

$$\{\mathbf{y}(t_k; \mathbf{y}_i) \mid t_k \in [t_i, t_{i+1}], \mathbf{y}_i \in [\mathbf{y}_i]\}$$

Phase 2 tight enclosure $[\mathbf{y}_{t_{i+1}}]$ at time t_{i+1} .

A trajectory then consists in a set of boxes called a **tube**.



If we consider again the dynamic of a drone i :

$$(\mathcal{S}_i) \begin{cases} \dot{\mathbf{x}}_i = f(\mathbf{x}_i, \mathbf{u}_i) \\ \mathbf{x}_i(0) \in [\mathbf{x}_{i,0}] \\ \mathbf{u}_i \in [\mathbf{u}_i] \end{cases}$$

The control \mathbf{u}_i is considered **constant** during the simulation over time t .

we infer a solution operator

$$\mathbf{x}_i(t, [\mathbf{x}_{i,0}], [\mathbf{u}_i]) = \{\mathbf{x}(t; \mathbf{x}_0; \mathbf{u}_i) \mid \mathbf{x}_0 \in [\mathbf{x}_{i,0}], \mathbf{u}_i \in [\mathbf{u}_i]\}.$$

and its associated interval inclusion $[\mathbf{x}_i](t, [\mathbf{x}_{i,0}], [\mathbf{u}_i])$.

Validated Simulation with Runge-Kutta

- Proof of existence and uniqueness of solution for ODEs and DAEs,
- Local truncation error computation for any Runge-Kutta method (implicit or explicit),
- Combined with *contractors* (HC4, etc.).

Verification of temporal constraints

- Stayed in \mathcal{A} until $\tilde{t} < t_{\text{end}}$:

$$\forall t \in [0, \tilde{t}], \{y(t; y_0) \mid y_0 \in [y_0]\} \subseteq \text{int}(\mathcal{A})$$

- Included in \mathcal{A} inside $[0, t_{\text{end}}]$:

$$\exists t \in [0, t_{\text{end}}], \{y(t; y_0) \mid y_0 \in [y_0]\} \subseteq \text{int}(\mathcal{A}).$$

Temporal constraint

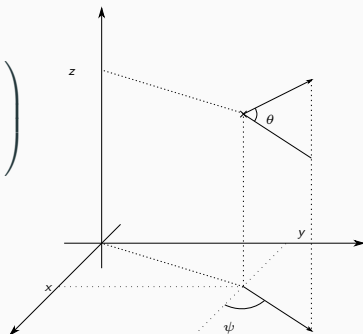
Has crossed \mathcal{A} (before τ):

$$\exists t < \tau, \mathbf{y}(t) \cap \mathcal{A} \neq \emptyset$$

We can define the collision detection in term of this temporal constraint.

Solving the Problem with DynIBEX: a Small Example

$$(S_i) \begin{cases} \dot{\mathbf{X}}_i = \begin{pmatrix} \dot{x}_i \\ \dot{y}_i \\ \dot{z}_i \end{pmatrix} = \begin{pmatrix} v_i \cos \psi_i \cos \theta_i \\ v_i \sin \psi_i \cos \theta_i \\ v_i \sin \theta_i \end{pmatrix} \\ \mathbf{X}_i(0) \in [\mathbf{X}_0] \end{cases}$$



with

- the state vector $\mathbf{X}_i = (x_i, y_i, z_i)^T$ representing the position of the robot;
- the control vector $\mathbf{u}_i = (v_i, \psi_i, \theta_i)^T$ consisting in the velocity v_i , the heading angle ψ_i and the track angle θ_i .

Solving the Problem with DynIBEX: a Small Example

drone (\mathcal{S}_1):

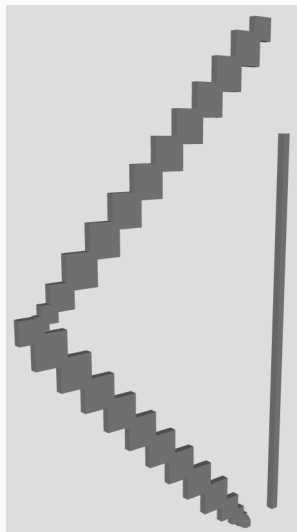
- $\mathbf{X}_1 = (1, 1, 1)$ and $v_1 = 1$;
- $\psi_1 = \frac{\pi}{2}$ and $\theta_1 = \frac{\pi}{2}$.

drone (\mathcal{S}_2):

- $\mathbf{X}_2 = (1, 7.8, 7.8)$ and $v_2 = -1$;
- $\psi_2 = -\frac{\pi}{2}$ and $\theta_2 = -\frac{3\pi}{4}$.

drone (\mathcal{S}_3):

- $\mathbf{X}_3 = (0, 1, 2)$ and $v_3 = 1$;
- $\psi_3 = \pi$ and $\theta_3 = \frac{\pi}{2}$.



The simulation time is 10s.

For N drones, we obtain the N tubes :

- $(\mathcal{S}_1) : \left\{ \left[\mathbf{x}_{1;[t_{1,0}]} \right], \left[\mathbf{x}_{1;[t_{1,1}]} \right], \dots, \left[\mathbf{x}_{1;[t_{1,m_1}]} \right] \right\}$
- $(\mathcal{S}_2) : \left\{ \left[\mathbf{x}_{2;[t_{2,0}]} \right], \left[\mathbf{x}_{2;[t_{2,1}]} \right], \dots, \left[\mathbf{x}_{2;[t_{2,m_2}]} \right] \right\}$
- \vdots
- $(\mathcal{S}_N) : \left\{ \left[\mathbf{x}_{N;[t_{N,0}]} \right], \left[\mathbf{x}_{N;[t_{N,1}]} \right], \dots, \left[\mathbf{x}_{N;[t_{N,m_N}]} \right] \right\}$

with $\left[\mathbf{x}_{i;[t_{i,j}]} \right] \supseteq \{ \mathbf{x}_i(t) \mid t \in [t_{i,j}] \}$

Algorithm 1: Algorithm for checking the non collision between a set of drones.

Input: $(S_1) : \left\{ \left[\mathbf{x}_{1;[t_{1,0}]} \right], \left[\mathbf{x}_{1;[t_{1,1}]} \right], \dots, \left[\mathbf{x}_{1;[t_{1,m_1}]} \right] \right\}$

Input: $(S_2) : \left\{ \left[\mathbf{x}_{2;[t_{2,0}]} \right], \left[\mathbf{x}_{2;[t_{2,1}]} \right], \dots, \left[\mathbf{x}_{2;[t_{2,m_2}]} \right] \right\}$

\vdots

Input: $(S_N) : \left\{ \left[\mathbf{x}_{N;[t_{N,0}]} \right], \left[\mathbf{x}_{N;[t_{N,1}]} \right], \dots, \left[\mathbf{x}_{N;[t_{N,m_N}]} \right] \right\}$

for $i = 1$ **to** N **do**

for $j = 0$ **to** m_i **do**

for $k = i + 1$ **to** N **do**

for $l = 0$ **to** m_k **do**

if $[t_{i;j}] \cap [t_{k;l}] \neq \emptyset$ **then**

if $\left[\mathbf{x}_{i;[t_{i;j}]} \right] \cap \left[\mathbf{x}_{k;[t_{k;l}]} \right] \neq \emptyset$ **then**

 possible collision

no collision detected

We end up with a solution with a **high complexity**.

DynIbex

- handle the dynamics;
- handle static obstacles constraints easily (for example boxRRT);
- requires more sophisticated tools to handle the dynamic obstacles.

A solution:

AbSolute

(Constraint solver based on abstract domains)

We cast the results from DynIBEX into a Constraint Satisfaction Problem.

init

```
real x0 = [-10000.000000; 10000.000000]
real x1 = [-10000.000000; 10000.000000]
real x2 = [-10000.000000; 10000.000000]
real t = [0; 100.0]
```

constraints

$T_0 :$

$(t \text{ in } [t_{0,0}] \ \&\& \ x_0 \text{ in } [x_{1:[t_{0,0}}]_0} \ \&\& \ x_1 \text{ in } [x_{1:[t_{0,0}}]_1} \ \&\& \ x_2 \text{ in } [x_{1:[t_{0,0}}]_2}) \parallel$

$(t \text{ in } [t_{0,1}] \ \&\& \ x_0 \text{ in } [x_{1:[t_{0,1}}]_0} \ \&\& \ x_1 \text{ in } [x_{1:[t_{0,1}}]_1} \ \&\& \ x_2 \text{ in } [x_{1:[t_{0,1}}]_2}) \parallel$

\vdots

$T_1 :$

$(t \text{ in } [t_{1,0}] \ \&\& \ x_0 \text{ in } [x_{1:[t_{1,0}}]_0} \ \&\& \ x_1 \text{ in } [x_{1:[t_{1,0}}]_1} \ \&\& \ x_2 \text{ in } [x_{1:[t_{1,0}}]_2}) \parallel$